

Query enumeration over Nowhere-Dense databases

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Join work with: Arnaud Durand, Nicole Schweikardt, Luc Segoufin

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Introduction

- Query q
- Database D
- Compute $q(D)$

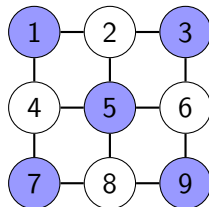
small
huge
gigantic

Examples :

query q

$$q(x, y) := \exists z(B(x) \wedge E(x, z) \wedge \neg E(y, z))$$

database D



solutions $q(D)$

$\{(1,2) (1,3) (1,4)$
 $(1,6) (1,7) \dots$
 $(3,1) (3,2) (3,4)$
 $(3,6) (3,7) \dots$
 $\dots \}$

Enumeration

Input : $\|D\| := n$ & $\|q\| := k$ (computation with RAM)

Goal : output solutions one by one (no repetition)

- STEP 1: Preprocessing

Prepare the enumeration : Database $D \longrightarrow$ Index I

Preprocessing time : $f(k) \cdot n \rightsquigarrow O(n)$

- STEP 2 : Enumeration

Enumerate the solutions : Index $I \longrightarrow \bar{x}_1, \bar{x}_2, \bar{x}_3, \bar{x}_4, \dots$

Delay : $O(f(k)) \rightsquigarrow O(1)$

Constant delay enumeration after linear preprocessing

Example

Input :

- Database $D := \langle \{1, \dots, n\}; E \rangle$ $\|D\| = |E|$ ($E \subseteq D \times D$)
- Query $q(x, y) := \neg E(x, y)$

D

(1,1)

(1,2)

(1,6)

⋮

(2,3)

⋮

(i,j)

(i,j+1)

(i,j+3)

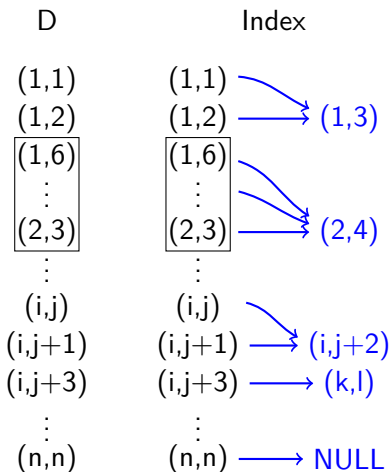
⋮

(n,n)

Example

Input :

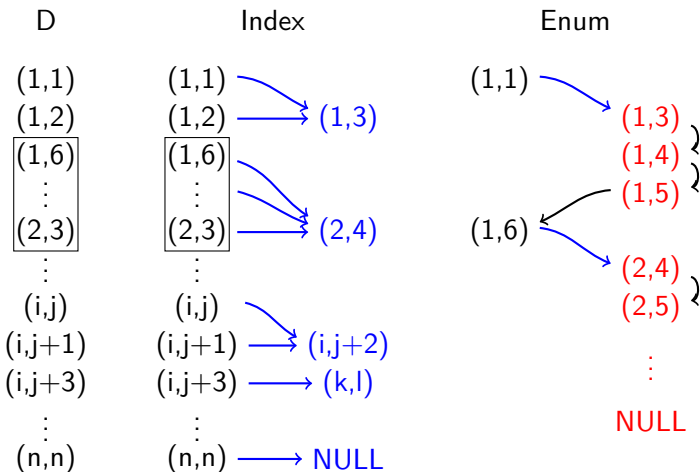
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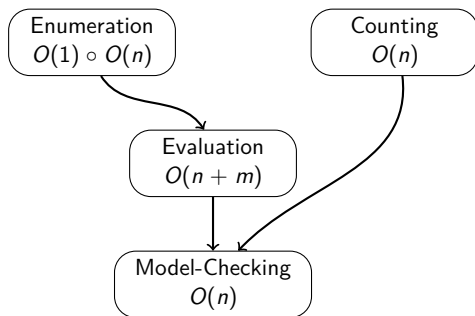
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Other problems

For **FO** queries over a class \mathcal{C} of databases.

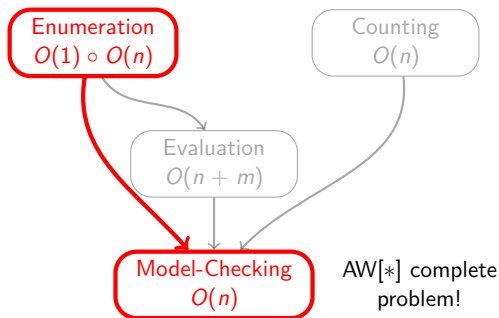
Model-Checking	: Is this true ?	$O(n)$
Enumeration	: Enumerate the solutions	$O(1) \circ O(n)$
Counting	: How many solutions ?	$O(n)$
Evaluation	: Compute the entire set	$O(n + m)$



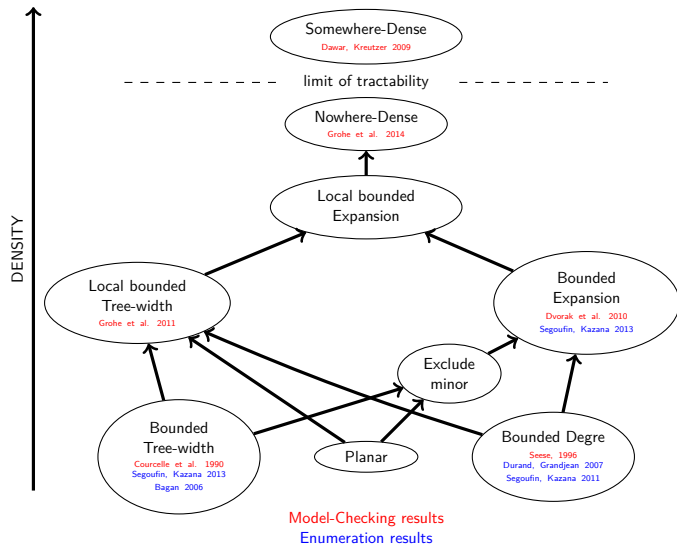
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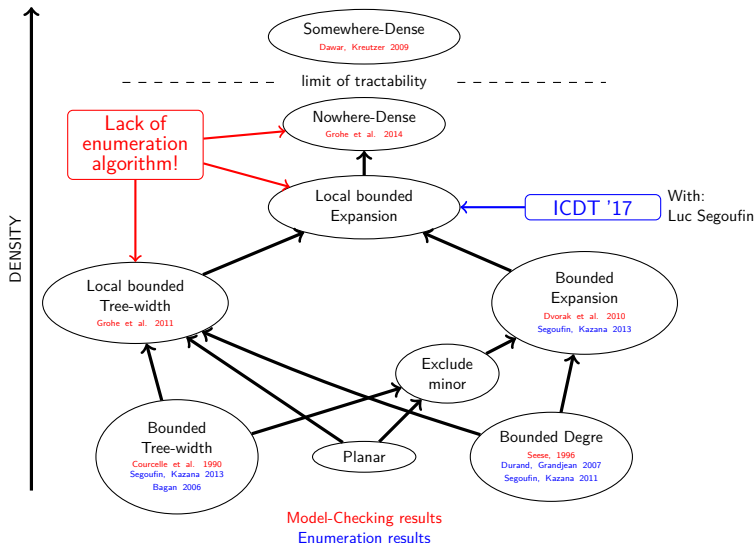
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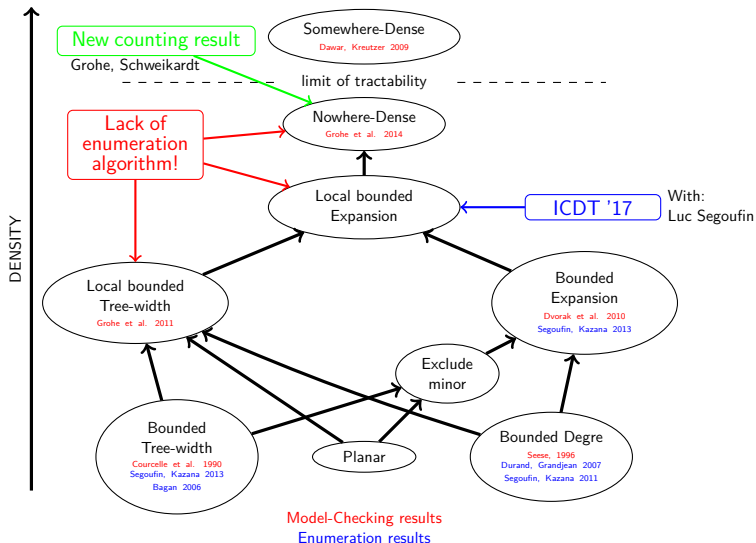
Classes of graphs closed under taking sub-graphs



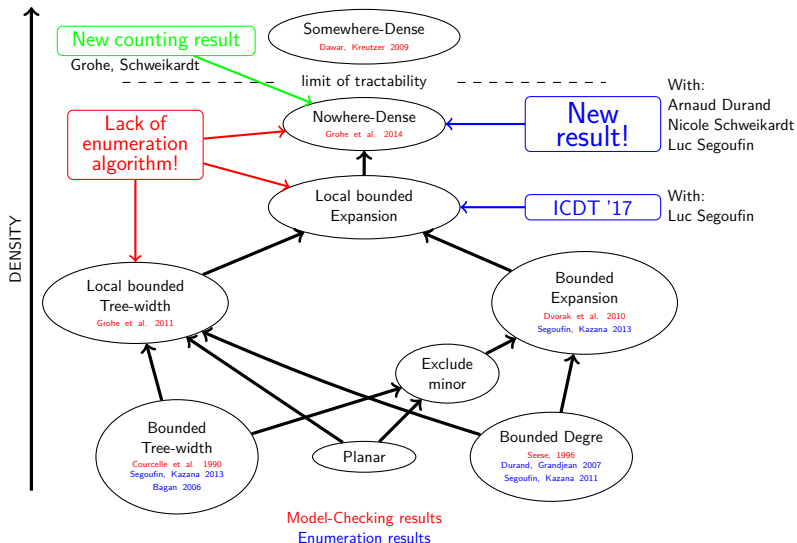
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Classes of graphs closed under taking sub-graphs



Tools and results (unpublished yet)

We use :

- A new “Hanf like” normal form for FO queries.¹
- The algorithm for the model checking.²
- Neighbourhood cover.²
- Game characterization of Nowhere-Dense classes.²
- Short-cut pointers dedicated to the enumeration.³

We can :

- Enumerate with constant delay after pseudo-linear preprocessing.
- Test in constant time after pseudo-linear preprocessing.

¹Grohe, Schweikardt '17

²Grohe, Kreutzer, Siebertz '14

³Segoufin, V. '17

Future work

Enumeration with update:

What happens if a small change occurs after the preprocessing ?

Existing results for: words, graphs with bounded tree-width or bounded degree.

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Thank you !

Questions ?